

QUESTION

- (a) State the Duality Theorem of linear programming and use it to prove the Theorem of Complementary Slackness.
- (b) Use duality theory to determine whether $x_1 = 0, x_2 = 1, x_3 = 0, x_4 = 4$, is an optimal solution of the linear programming problem

$$\begin{aligned} \text{maximize} \quad & z = 4x_1 + x_2 + 7x_3 + 9x_4 \\ \text{subject to} \quad & x_1 \geq 0, x_2 \geq 0, x_3 \geq 0, x_4 \geq 0 \\ & 6x_1 + 8x_2 + 3x_3 + x_4 \leq 15 \\ & 3x_1 + 2x_2 + 7x_3 + 4x_4 \leq 18 \\ & 5x_1 + 5x_2 + 8x_3 + 3x_4 = 17. \end{aligned}$$

- (c) For the linear programming problem

$$\begin{aligned} \text{maximize} \quad & \sum_{j=1}^n c_j x_j \\ \text{subject to} \quad & x_j \geq 0 \quad j = 1, \dots, n \\ & \sum_{j=1}^n a_{ij} x_j \leq b_i \quad i = 1, \dots, m, \end{aligned}$$

the optimal value of the objective function is z^* and y_1^*, \dots, y_m^* are optimal values of the dual variables. Let z^{**} denote the optimal value of the objective function for the linear programming problem

$$\begin{aligned} \text{maximize} \quad & \sum_{j=1}^n c_j x_j \\ \text{subject to} \quad & x_j \geq 0 \quad j = 1, \dots, n \\ & \sum_{j=1}^n a_{ij} x_j \leq b_i + \delta_i \quad i = 1, \dots, m. \end{aligned}$$

$$z^{**} \leq z^* + \sum_{i=1}^m \delta_i y_i^*.$$

You may use the Duality Theorem in your proof.

ANSWER

- (a) The duality theorem states that

- if the primal problem has an optimal solution, then so has the dual, and $z_p = z_D$;
- if the primal problem is unbounded, then the dual is infeasible;
- if the primal problem is infeasible, then the dual is either infeasible or unbounded.

Consider the following primal and dual problems

$$\begin{array}{ll}
\text{Maximize} & z_P = \mathbf{c}^T \mathbf{x} \\
\text{subject to} & \mathbf{x} \geq \mathbf{0}, \mathbf{s} \geq \mathbf{0} \\
& A\mathbf{x} + \mathbf{s} = \mathbf{b}
\end{array}
\quad
\begin{array}{ll}
\text{Minimize} & z_D = \mathbf{b}^T \mathbf{y} \\
\text{subject to} & \mathbf{y} \geq \mathbf{0}, \mathbf{t} \geq \mathbf{0} \\
& A^T \mathbf{y} - \mathbf{t} = \mathbf{c}
\end{array}$$

For feasible solutions of the primal and dual, we have

$$Z_P = \mathbf{c}^T \mathbf{x} = (\mathbf{y}^T A - \mathbf{t}^T) \mathbf{x} = \mathbf{y}^T (\mathbf{b} - \mathbf{s}) - \mathbf{t}^T \mathbf{x} = z_D - \mathbf{y}^T \mathbf{s} - \mathbf{t}^T \mathbf{x}$$

For an optimal solution of the primal and dual, $z_p = z_D$ so

$$\mathbf{y}^T \mathbf{s} + \mathbf{t}^T \mathbf{x} = 0$$

Since variables are non negative this implies that

$$y_i s_i = 0 \quad i = 1, \dots, m$$

$$t_j x_j = 0 \quad j = 1, \dots, n$$

- (b) The solution $x_1 = 0, x_2 = 1, x_3 = 0, x_4 = 4$ yeilds $z = 37, s_1 = 3, s_2 = 0$.

The dual problem is

$$\begin{array}{ll}
\text{minimize} & z_D = 15y_1 + 18y_2 + 17y_3 \\
\text{subject to} & y_1 \geq 0, y_2 \geq 0, \\
& 6y_1 + 3y_2 + 5y_3 \geq 4 \\
& 8y_1 + 3y_2 + 5y_3 \geq 1 \\
& 3y_1 + 7y_2 + 8y_3 \geq 7 \\
& y_1 + 4y_2 + 3y_3 \geq 9
\end{array}$$

If the given solution is optimal, then we can use the complementary slackness conditions.

$$y_1 s_1 = 0 \text{ implies } y_1 = 0$$

$$x_2 t_2 = 0 \text{ implies } t_2 = 0$$

$$x_4 t_4 = 0 \text{ implies } t_4 = 0$$

Thus,

$$\begin{array}{rcl}
2y_2 + 5y_3 & = & 1 \\
4y_2 + 3y_3 & = & 9
\end{array}$$

so

$$y_2 = 3, y_3 = -1$$

and

$$t_1 = 0, t_3 = 6$$

Therefore, the solution is feasible.

Since $z_D = 37 = z$ the proposed solution is optimal.

(c) Using matrix notation, the relevant problems are

$$\begin{array}{ll} \text{(P)} & \text{Maximize } \mathbf{c}^T \mathbf{x} \\ & \text{subject to } \mathbf{x} \geq \mathbf{0} \\ & \quad \quad \quad \mathbf{Ax} \leq \mathbf{b} \end{array} \quad \begin{array}{ll} \text{(P')} & \text{Maximize } \mathbf{c}^T \mathbf{x} \\ & \text{subject to } \mathbf{x} \geq \mathbf{0} \\ & \quad \quad \quad \mathbf{Ax} \leq \mathbf{b} + \delta \end{array}$$

and the dual of (P) is

$$\begin{array}{ll} \text{(D)} & \text{Minimize } \mathbf{b}^T \mathbf{y} \\ & \text{subject to } \mathbf{y} \geq \mathbf{0} \\ & \quad \quad \quad \mathbf{A}^T \mathbf{y} \geq \mathbf{c} \end{array}$$

From the duality theorem,

$$z^* = \mathbf{b}^T \mathbf{y}^* = (\mathbf{y}^*)^T \mathbf{b}$$

Let $\mathbf{x} = \mathbf{x}^*$ be an optimal solution of (P'). Then

$$x^{**} = \mathbf{c}^T \mathbf{x}^* \leq (\mathbf{y}^*)^T \mathbf{Ax}^* \leq (\mathbf{y}^*)^T (\mathbf{b} + \delta) = z^* + \sum_{i=1}^m s_i y_i^*$$